

WISST! The Lefschetz principle

or: What is... quantifier elimination?

Given polynomials f_1, \dots, f_r and $g_1, \dots, g_s \in \mathbb{Z}[x_1, \dots, x_n]$
study the solvability of the system

$$\begin{aligned} f_i &= 0, \quad i \in [r], \\ g_j &\neq 0, \quad j \in [s], \end{aligned} \quad (\exists)$$

over varying fields \mathbb{K} . **algebraically closed** fields \mathbb{K} .

Note: the \mathbb{Z} coefficients make sense in every field \mathbb{K} by interpreting $k \in \mathbb{Z}$ as $\underbrace{1 + 1 + \dots + 1}_{k \text{ times}} \in \mathbb{K}$.

The Lefschetz principle

1. The answer to (\exists) depends only on the characteristic of the field. If a solution exists over \mathbb{K} , then also over $\overline{\mathbb{K}}$ (i.e. $\overline{\mathbb{Q}}$ or $\overline{\mathbb{F}_p}$).
2. If (\exists) has a solution in characteristic 0, then it has a solution over characteristic $p > 0$ for all but finitely many primes p .
3. In particular if (\exists) has a solution over \mathbb{C} , it has a solution in finite fields \mathbb{F}_{p^m} for all but finitely many primes p .
4. The set of characteristics over which (\exists) has a solution can be effectively computed from the defining polynomials f_i and g_j .

Polynomial systems and Model theory

Model theory studies “models”: mathematical objects **interpreting** symbols of a formal language and **satisfying** axioms written in that language as well as the **theory** of a given model, i.e. all formulas in the language which are true in this model.

(First-order) language of rings:

- ▶ constants 0 and 1
- ▶ functions $+$, $-$, \cdot
- ▶ relations $=$
- ▶ Boolean logic connectives \wedge , \vee , \neg (\Rightarrow , \Leftrightarrow , \dots)
- ▶ \exists and \forall quantifiers
- ▶ variables x_1, x_2, \dots

Formulas, axioms and definability

The definition of a ring can be written in the language of rings:

- ▶ $\forall a : a - a = 0$
- ▶ $\forall a, b, c : (a + b) + c = a + (b + c)$
- ▶ $\forall a, b : a + b = b + a$
- ▶ ...

A field is a commutative ring with inverses:

- ▶ $\forall a \exists b : \neg(a = 0) \Rightarrow ab = 1$

(\exists) is expressible as a sentence in this language:

$$\exists x_1, \dots, x_n : \bigwedge_{i=1}^r f_i(x_1, \dots, x_n) = 0 \wedge \bigwedge_{j=1}^s \neg(g_j(x_1, \dots, x_n) = 0)$$

Quantifier elimination

The theory of *algebraically closed* fields admits **quantifier elimination** in the language of rings. That is, for each formula φ there exists an equivalent formula ψ without \exists or \forall quantifiers, such that for every algebraically closed field \mathbb{K} :

$$\mathbb{K} \models \varphi \Leftrightarrow \mathbb{K} \models \psi.$$

Eliminating quantifiers (and hence **all variables**) from (\exists) results in a Boolean combination of (in)equations $n = m$ for some $n, m \in \mathbb{Z}$.

These inequalities point out exactly which **characteristics** are required and which are ruled out for having a solution to (\exists) .