

The Ingleton inequality for random variables

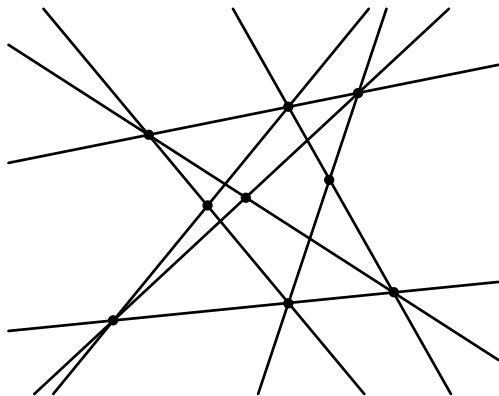
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13 August 2024

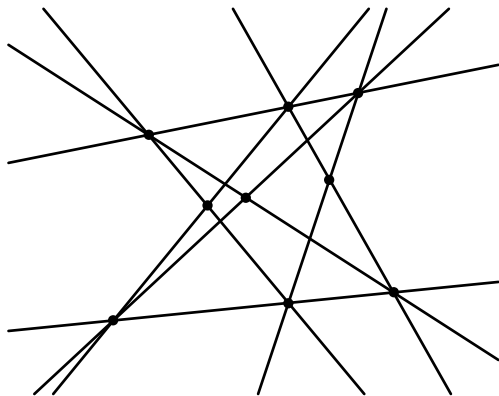
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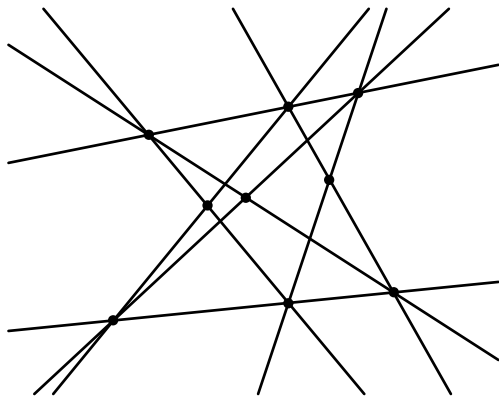
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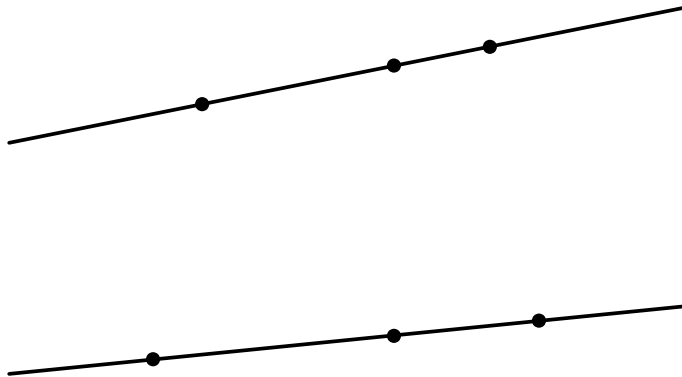


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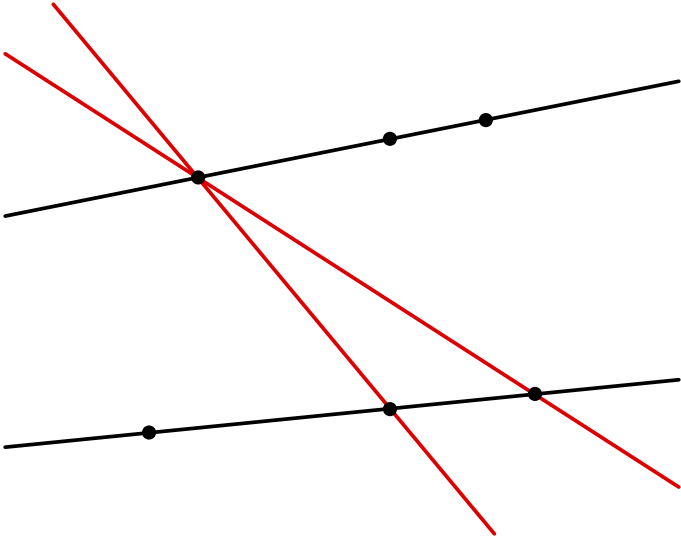
- ▶ Matroids are combinatorial structures which model “special position” relations in geometry.
 - ▶ For example the matroid of a set of points in the projective plane records which triples of points lie on a line.
- ▶ Non-realizability of matroids captures the (non-obvious) laws of projective geometry.



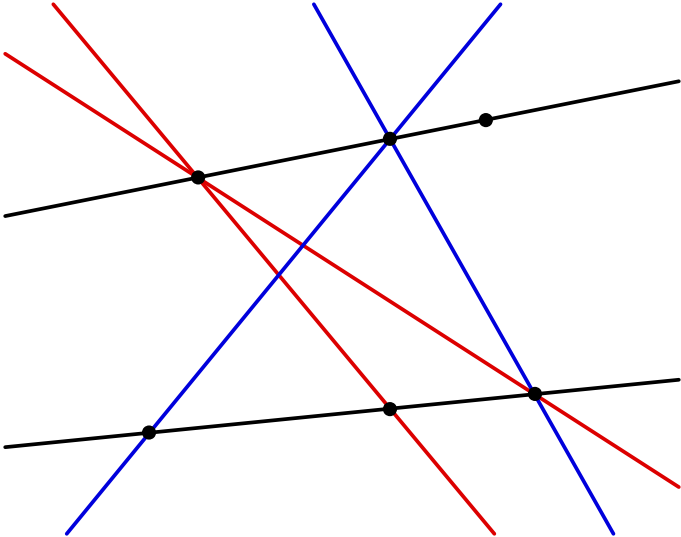
Laws of geometry



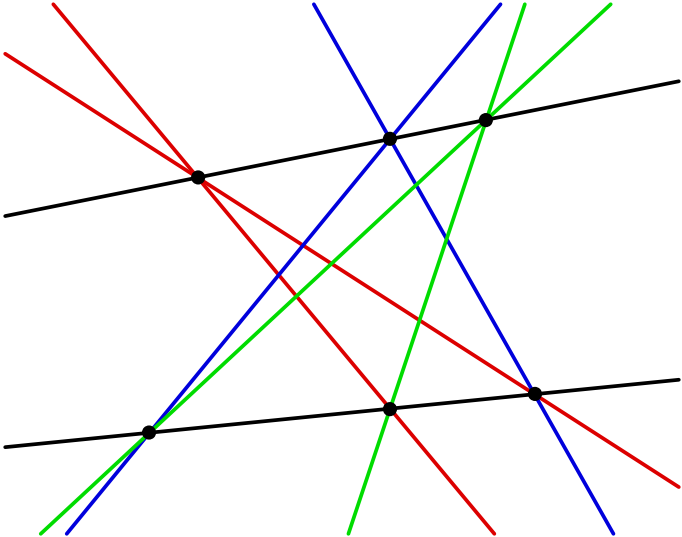
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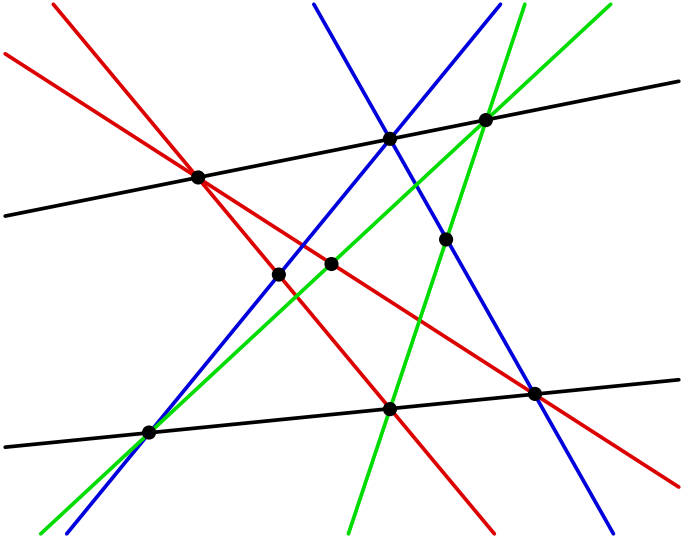
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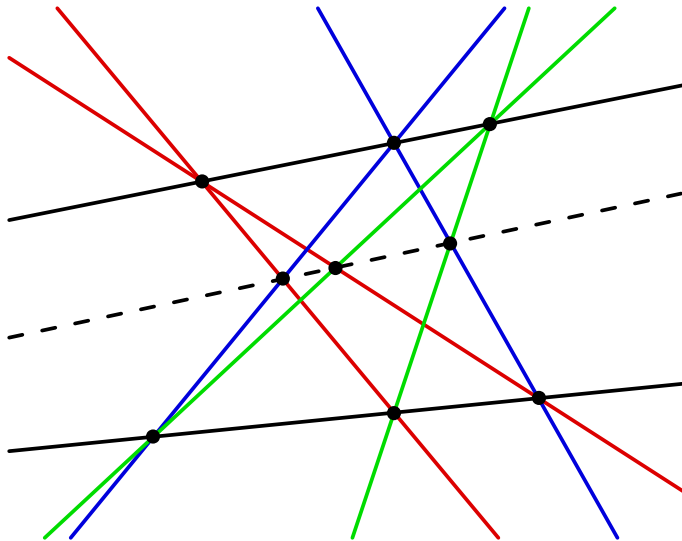
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Entropy

Let X be a random variable taking finitely many values $\{1, \dots, d\}$ with positive probabilities. Its *Shannon entropy* is

$$H(X) := \sum_{i=1}^d p(X = i) \log 1/p(X = i).$$

- ▶ H is continuous on $\Delta(d)$ and analytic on the interior.

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- ▶ H is continuous on $\Delta(d)$ and analytic on the interior.
- ▶ A random vector $X \in \Delta(d_1, \dots, d_n)$ is a random variable in $\Delta(\prod_{i=1}^n d_i)$, so the definition of H extends to vectors.
- ▶ The random vector $X = (X_1, \dots, X_n)$ has 2^n marginal random vectors and we collect their entropies in an **entropy profile** $h_X : 2^{[n]} \rightarrow \mathbb{R}$.
 - ▶ For example (X, Y) has entropy profile $(0, H(X), H(Y), H(X, Y)) \in \mathbb{R}^4$.

Entropy as information

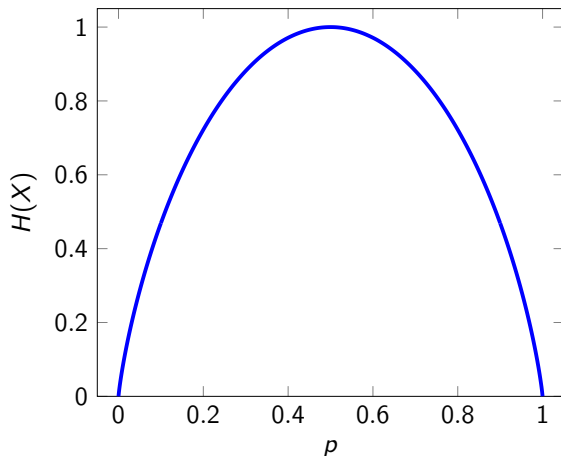


Figure: Entropy of a binary random variable X as a function of $p = p(X = \text{heads})$.

The entropy region and information inequalities

Let $\mathbf{H}_n^* \subseteq \mathbb{R}^{2^n}$ consist of all h_X where X is an n -variate discrete random vector. \mathbf{H}_n^* is the image of $\bigcup_{d_1=1}^{\infty} \cdots \bigcup_{d_n=1}^{\infty} \Delta(d_1, \dots, d_n)$ under the transcendental map $X \mapsto h_X$.

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- ▶ Elements of the dual cone ([linear information inequalities](#)) can give bounds for optimization problems.

Shannon inequalities

- ▶ A function $h : 2^N \rightarrow \mathbb{R}$ is a **polymatroid** if
 - ▶ $h(\emptyset) = 0$,
 - ▶ $h(I | K) := h(IK) - h(K) \geq 0$ for disjoint I and K ,
 - ▶ $h(I : J | K) := h(IK) + h(JK) - h(IJK) - h(K) \geq 0$ for disjoint I, J, K .

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- ▶ The set \mathbf{P}_N of polymatroids is a polyhedral cone in \mathbb{R}^{2^N} and $\mathbf{P}_N \supseteq \overline{\mathbf{H}_N^*} \rightarrow \text{ITIP}$.
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Theorem ([Mat07])

$\overline{\mathbf{H}_N^*}$ is not polyhedral for $|N| \geq 4$.

- ▶ **GMM conjecture**: $\overline{\mathbf{H}_N^*}$ is not semialgebraic for $|N| \geq 4$.

Information inequalities abound

Rule [43] Given:

$$\begin{aligned} & aI(A; B) \\ \leq & bI(A; B|C) + cI(A; C|B) + zI(B; C|A) \\ & + eI(A; B|D) + fI(A; D|B) \\ & + (b' + d' + z)I(B; D|A) + hI(C; D) \\ & + iI(C; D|A) + zI(C; D|B) \end{aligned}$$

and

$$\begin{aligned} & a'I(A; B) \\ \leq & b'I(A; B|C) + c'I(A; C|B) + d'I(B; C|A) \\ & + e'I(A; B|D) + f'I(A; D|B) + g'I(B; D|A) \\ & + h'I(C; D) + i'I(C; D|A) + j'I(C; D|B) \end{aligned}$$

Get:

$$\begin{aligned} & (a + a' + z)I(A; B) \\ \leq & (a + b + c + f + b' + 2z)I(A; B|C) \\ & + (-a + b + c + e + e' + z)I(A; C|B) \\ & + (d' + z)I(B; C|A) + (e + e' + z)I(A; B|D) \\ & + (f + f')I(A; D|B) \\ & + (-a' + b' + e' + g' + i')I(B; D|A) \\ & + (h + h' + z)I(C; D) + (i + i')I(C; D|A) \\ & + (j')I(C; D|B) \end{aligned}$$

Using: RS is copy of CD over AB

Substitutions: $A C R S; AD B R S$

Abbreviated Proof of (75): T : D-copy of A over BCRS.

L1: $-a.c. +c.d. +r.cd.a +c.s.a +b.d.s +a.bs.d +2a.cr.bs +a.bs.cr +d.r.abes +d.s.abcr$

SL1: $d.t.a +c.d.t +a.t.cd +c.r.t +a.t.cr +d.r.act +b.t.acdr +a.t.bs +c.s.at +b.t.aes +d.t.s +a.s.dt +b.d.ast +c.t.abds +a.r.best +r.ad.best +s.ad.bert +d.t.abers$ C2L1: $3t.ad.bcrs$

S: C-copy of A over BDR.

L2: $-2a.c. +2c.d. +a.b.cr +2a.c.br +c.ar.b +a.b.dr +4a.d.br +2a.br.d +2d.br.a +2r.cd.a +d.r.abc$

SL2: $c.s.b +a.b.cs +c.d.s +a.s.cd +d.s.abc +3a.s.br +3c.s.br +c.r.abs +d.r.s +a.s.dr +d.r.abs +d.br.as +c.r.ads +b.s.acdr +2c.s.abdr +2d.s.abcr$

C2L2: $7s.ac.bdr$

R: D-copy of C over AB .

S: $c.r.a +3c.r.b +d.r.a +7d.r.b +c.d.r +2b.r.acd +r.ab.cd +9c.r.abd +3d.r.abc$

C2: $16r.cd.ab$

Randall Dougherty, Chris Freiling, and Kenneth Zeger. *Non-Shannon Information Inequalities in Four Random Variables*. 2011. arXiv: 1104.3602v1 [cs.IT]

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All of these are **linear** in h . Even though entropy is a transcendental function, many of these conditions are **polynomial** in the probabilities \rightarrow algebraic statistics.

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- ▶ Let A, B, C, D be subspaces in a finite-dimensional vector space.
Then the **Ingleton inequality** holds for $h = \dim$ (the matroid setting):

$$\begin{aligned} I(A, B \mid C, D) := & h(A, C) + h(B, C) + h(A, D) + h(B, D) + h(C, D) - \\ & h(A, B) - h(C) - h(D) - h(A, C, D) - h(B, C, D) \geq 0. \end{aligned}$$

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These are **conditional linear information inequalities** and they can sometimes tell apart honest boundary parts of \mathbf{H}_n^* from fake boundary parts on $\overline{\mathbf{H}_n^*}$.

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 - ▶ Six linearly isomorphic copies of the cone $\mathbf{J}_4 := \mathbf{P}_4 \cap \{I(A, B \mid C, D) < 0\}$.
- ▶ \mathbf{J}_4 is simplicial and its facets are induced by **conditional independence** functionals and the **Ingleton** functional.

Conditional Ingleton inequalities

Theorem ([KR13] & [Stu21] & [Boe23])

Up to symmetry there are precisely ten minimal sets of conditional independence assumptions on four random variables which ensure $I \geq 0$.

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Problem

Which of these laws holds on $\overline{\mathbf{H}}_4^$? (Some do, some don't ...)*

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Find/sample positive points from conditional independence varieties.

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Is $[A \perp\!\!\!\perp C \mid D] \wedge [A \perp\!\!\!\perp D \mid C] \wedge [B \perp\!\!\!\perp C \mid D] \wedge [B \perp\!\!\!\perp D \mid C] \implies I(A, B \mid C, D) \geq 0$ essentially conditional?

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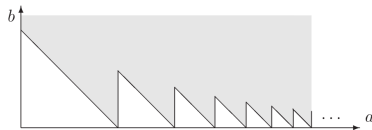
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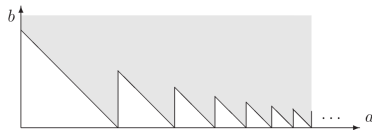
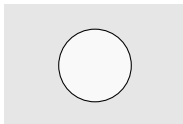
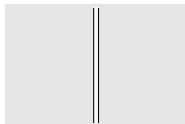
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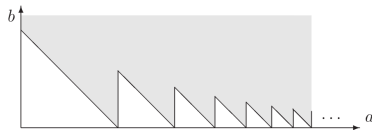
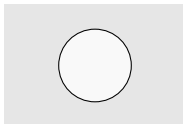
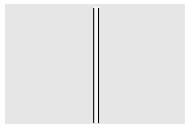
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Thank you!

References

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