

Matroids and chromatic polynomials

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Geometry of Lorentzian polynomials day
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Combinatorial geometries

- ▶ Matroids arise as a common combinatorial abstraction of **independence** in linear algebra and graph theory.
- ▶ But appear also in numerous other contexts such as algebraic or stochastic independence, game theory, rigidity theory, cryptography ...

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- ▶ Example: Let (v_1, \dots, v_n) be elements of a vector space. The subsets $I \subseteq [n]$ such that $(v_i : i \in I)$ are linearly independent form a matroid.
- ▶ Example: Let $G = (V, E)$ be an undirected graph. The sets $I \subseteq E$ which do not contain a cycle form a matroid.

Matroid cryptomorphisms I: Independent sets

A matroid on ground set N is specified by its collection of **independent sets** $\mathcal{J} \subseteq 2^N$ satisfying:

- ▶ $\emptyset \in \mathcal{J}$,
- ▶ $J \subseteq I \in \mathcal{J}$ implies $J \in \mathcal{J}$,
- ▶ if $I, J \in \mathcal{J}$ with $|I| > |J|$, then there exists $x \in I \setminus J$ such that $J \cup \{x\} \in \mathcal{J}$.

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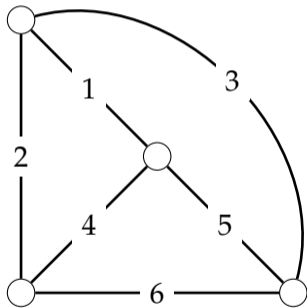
The independent sets of matroids are special **(pure) simplicial complexes**:

n	1	2	3	4	5	6
Simplicial complexes	1	2	5	20	180	16 143
Matroids	1	2	4	9	21	60

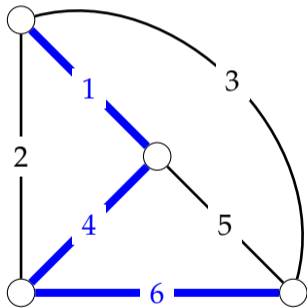
Nelson (2016): Almost all matroids are not linearly representable.

Example

Independent sets:



Example



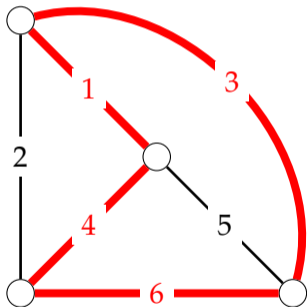
Independent sets:

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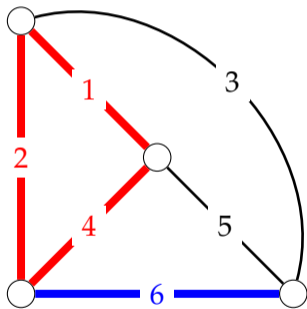
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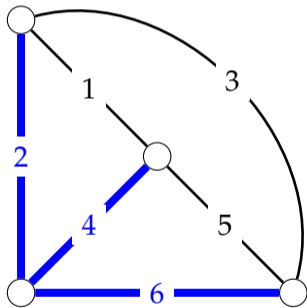
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Example



Independent sets:

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- ▶ $\{2,4,6\} \dots$

Matroid cryptomorphisms II: Rank function

A matroid on ground set N is specified by its **rank function** $r : 2^N \rightarrow \mathbb{Z}$ satisfying:

- ▶ $r(\emptyset) = 0$,
- ▶ $r(A) \leq r(B)$ for $A \subseteq B$,
- ▶ $r(A) + r(B) \geq r(A \cup B) + r(A \cap B)$,
- ▶ $r(A) \leq |A|$.

Matroid cryptomorphisms II: Rank function

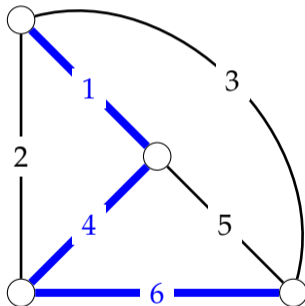
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Equivalence of independent sets and rank:

- ▶ A set I is independent if and only if $r(I) = |I|$.
- ▶ The rank of any set $A \subseteq N$ is $\max \{ |I| : A \supseteq I \in \mathcal{I} \}$.

Example



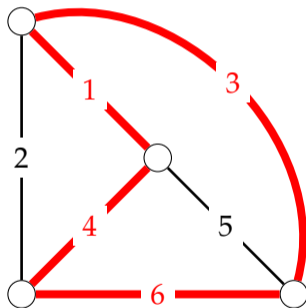
Independent sets:

- ▶ $\{1, 4, 6\}$
- ▶ $\{2, 4, 6\} \dots$

Ranks:

- ▶ $r(\{1, 4, 6\}) = 3$

Example



Independent sets:

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- ▶ $\{2, 4, 6\} \dots$

Ranks:

- ▶ $r(\{1, 4, 6\}) = 3$
- ▶ $r(\{1, 3, 4, 6\}) = 3 < 4 \dots$

Matroid cryptomorphisms III: Closure operator

A matroid on ground set N is specified by its **closure operator** $c : 2^N \rightarrow 2^N$ satisfying:

- ▶ $A \subseteq c(A)$,
- ▶ $c(A) = c(c(A))$,
- ▶ $c(A) \subseteq c(B)$ for $A \subseteq B$,
- ▶ if $x \in c(A \cup \{y\}) \setminus c(A)$ then $y \in c(A \cup \{x\}) \setminus c(A)$.

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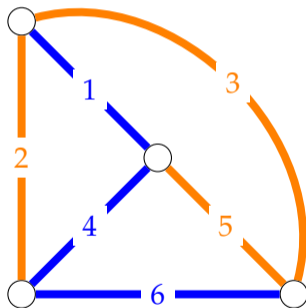
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Equivalence of rank and closure operator:

- ▶ A is closed if and only if $r(A \cup \{x\}) > r(A)$ for all $x \notin A$.
- ▶ The rank of any set $A \subseteq N$ is its rank in the lattice of flats.

Example



Independent sets:

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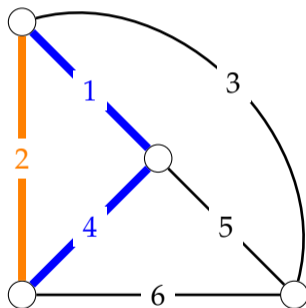
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- ▶ $r(\{1, 3, 4, 6\}) = 3 < 4 \dots$

Closures:

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Some terminology

- ▶ $x \in N$ is a **loop** if $r(\{x\}) = 0$.
- ▶ $x \neq y \in N$ are **parallel** if they are not loops and $r(\{x, y\}) = 1$.
- ▶ A matroid is **simple** if it has neither loops or parallel elements.



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Given a matroid M on N with rank function r and an element $x \in N$ we define two new matroids on $N \setminus \{x\}$:

- ▶ The **deletion** of x defines $M \setminus x$ whose rank function is the restriction $r|_{2^{N \setminus \{x\}}}$.
- ▶ The **contraction** of x defines M / x with rank function $r(A \cup \{x\}) - r(\{x\})$.

Chromatic polynomial of a graph

The **chromatic polynomial** of a graph $G = (V, E)$ is

$$\chi_G(q) := \# \text{ proper colorings of } G \text{ with } q \text{ colors.}$$

This is a polynomial because (provided x is neither a loop nor a coloop)

$$\begin{aligned}\chi_G(q) &= \chi_{G \setminus x}(q) - \chi_{G/x}(q), \\ \chi_G(q) &= q^n \text{ if } |V| = n \text{ and } E = \emptyset.\end{aligned}$$

Chromatic polynomial of a matroid


Chromatic (or [characteristic](#)) polynomial of a matroid: (See also: [Tutte polynomial](#).)


$$\begin{aligned}\chi_M(q) &= \chi_{M \setminus x}(q) - \chi_{M/x}(q) \\ &= \sum_{A \subseteq N} (-1)^{|A|} q^{r(N)-r(A)} =: \sum_{k=0}^{r(N)} w_k q^{r(N)-k}.\end{aligned}$$


w_k are the Whitney numbers of the first kind and they have alternating signs.

Teaser: Adiprasito-Huh-Katz: The sequence $|w_k|$ is log-concave and unimodal.

Further reading

 Karim Adiprasito, June Huh, and Eric Katz. “Hodge theory for combinatorial geometries”. English. In: *Ann. Math. (2)* 188.2 (2018), pp. 381–452. ISSN: 0003-486X. DOI: 10.4007/annals.2018.188.2.1.

 Federico Ardila. “The geometry of matroids”. English. In: *Notices Am. Math. Soc.* 65.8 (2018), pp. 902–908. ISSN: 0002-9920. DOI: 10.1090/noti1720.

 Dominic J. A. Welsh. *Matroid theory*. Vol. 8. London Mathematical Society Monographs. Academic Press, 1976.