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# Computational methods in Gaussian conditional independence inference

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MATHEMATISCHE  
KOMPLEXITÄTSREDUKTION

# Gaussian conditional independence

- Finite ground set  $[n] = \{1, \dots, n\}$  indexing some objects  $\xi_i$ .
- CI statement  $(ij|K)$  abbreviating  $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ .
- All CI statements  $\mathcal{A}_n = \{(ij|K) : ij \in \binom{[n]}{2}, K \subseteq [n] \setminus ij\}$ .

Let  $\xi \sim \mathcal{N}(\mu, \Sigma)$  be a vector of  $n$  random variables with a joint regular Gaussian distribution ( $\Sigma \in \text{PD}_n$ ), then

$$\begin{array}{ccc} \xi_i \perp\!\!\!\perp \xi_j \mid \xi_K & \Leftrightarrow & \det \Sigma_{iK, jK} = 0. \\ \Updownarrow & & \Downarrow \\ (ij|K) & & \det \Sigma_{ij|K} \end{array}$$

# Realizability and the Implication problem

$$[[\Sigma]] = \{(ij|K) \in \mathcal{A}_n : \det \Sigma_{ij|K} = 0\}.$$

The sets  $[[\Sigma]]$  for  $\Sigma \in \text{PD}_n$  have a special combinatorial structure.

For example:

$$(ij|L) \in [[\Sigma]] \wedge (ij|kL) \in [[\Sigma]] \Rightarrow (ik|L) \in [[\Sigma]] \vee (jk|L) \in [[\Sigma]] \quad (*)$$

$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L) \quad (*)$$

for all distinct  $i, j, k \in [n]$  and  $L \subseteq [n] \setminus ijk$ .

**Goal.** Describe all the sets  $[[\Sigma]]$ , the CI structures *realizable* by regular Gaussian distributions.

**Observation.** Knowing all realizable structures is equivalent to knowing all valid inference rules of the form  $(*)$ .

All valid inference rules for Gaussians can be “symmetrized”:

$$(12|) \wedge (12|3) \Rightarrow (13|) \vee (23|) \text{ for } n = 3 \curvearrowright$$
$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L) \quad \forall i, j, k, L \text{ for every } n \geq 3.$$

**Theorem (Šimeček / Sullivan).** No finite set of inference rules exactly describes regular Gaussian CI structures.

**Theorem.** Every finite set of inference rules which allows all regular Gaussian CI structures allows asymptotically  $2^{2^{\Omega(n)}}$  structures. The number of Gaussian CI structures is bounded by  $2^{\mathcal{O}(n^3)}$ .

**Theme:** properties of realizable structures give valid inference rules.  
Derive valid rules from approximations of realizable structures.

- 1 Combinatorial compatibility & SAT solvers
- 2 Information inequalities & LP solvers
- 3 The selfadhesivity phenomenon
- 4 Real algebra & Semidefinite programming

# Combinatorial compatibility

The *Gaussian CI configuration space* is the image of

$$\Sigma \mapsto (p_I(\Sigma) : I \subseteq [n]) \cup (a_{ij|K}(\Sigma) : (ij|K) \in \mathcal{A}_n).$$

Study the zero patterns of configuration vectors.

There exist polynomial relations on them, for example

$$p_{kL} a_{ij|L} - p_L a_{ij|kL} - a_{ik|L} a_{jk|L} = 0. \quad (\text{ET})$$

*Combinatorial compatibility* means “vanishing under uncertainty”:

What if we only knew that  $p_I \neq 0$  and whether  $a_{ij|K} = 0$  or  $a_{ij|K} \neq 0$ ?

$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L)$$

$$(ij|L) \wedge (ik|L) \Rightarrow (ij|kL)$$

$\vdots$

# Gaussoids

$$(ij|L) \wedge (ik|jL) \Rightarrow (ik|L) \wedge (ij|kL)$$

$$(ij|kL) \wedge (ik|jL) \Rightarrow (ij|L) \wedge (ik|L)$$

$$(ij|L) \wedge (ik|L) \Rightarrow (ij|kL) \wedge (ik|jL)$$

$$(ij|L) \wedge (ij|kL) \Rightarrow (ik|L) \vee (jk|L)$$

For  $n = 4$ :

- All subsets of  $\mathcal{A}_4$ : 16,777,216
- Gaussoids: 679
- Realizable gaussoids: 629

For  $n = 5$ :

- All subsets of  $\mathcal{A}_5$ : 1,208,925,819,614,629,174,706,176
- Gaussoids: 60,212,776
- Realizable gaussoids: ???

**Tools:** AllSAT solver to enumerate satisfying assignments.

# Breakout break



## Oriented gaussoids

$$p_{kL}a_{ij|L} - p_{L}a_{ij|kL} - a_{ik|L}a_{jk|L} = 0 \quad (\text{ET})$$

What if we only knew that  $\text{sgn } p_I = +1$  and the value of  $\text{sgn } a_{ij|K}$ ?

$$\begin{aligned} + (ij|L) \wedge - (ij|kL) &\Rightarrow [+(ik|L) \wedge +(jk|L)] \\ &\vee [-(ik|L) \wedge -(jk|L)] \end{aligned}$$

For  $n = 4$ :

- Gaussoids: 679
- Orientable gaussoids: **629**

For  $n = 5$ :

- Gaussoids: 60,212,776
- Orientable gaussoids: 20,584,290

**Tools:** AllSAT and incremental SAT solver to compute projected satisfying assignments.

# The multiinformation region

$$a_{ij|L}^2 = p_{iL}p_{jL} - p_{LPijL} \quad (\text{FT})$$

The *Gaussian multiinformation region* is the image of

$$\Sigma \mapsto (\log p_I(\Sigma) : I \subseteq [n]).$$

By (FT), the Gaussian multiinformation functions  $m(I) = \log \det \Sigma_I$  satisfy the following *linear information inequalities*

$$\Delta_{ij|K} m := m(iK) + m(jK) - m(ijK) - m(K) \geq 0, \quad (\text{Submodularity})$$

and submodularity at  $\Delta_{ij|K}$  is tight if and only if  $(ij|K) \in \llbracket \Sigma \rrbracket$ .

# Information inequalities

Linear information inequalities at the multiinformation region of the form

$$\sum_{\beta \in \mathcal{M}} c_{\beta} \Delta_{\beta} m \leq \sum_{\alpha \in \mathcal{L}} c_{\alpha} \Delta_{\alpha} m, \quad c_{\alpha}, c_{\beta} > 0, \quad \forall m$$

with  $\mathcal{L}, \mathcal{M} \subseteq \mathcal{A}_n$  encode CI inference rules

$$\bigwedge_{\alpha \in \mathcal{L}} \alpha \Rightarrow \bigwedge_{\beta \in \mathcal{M}} \beta.$$

CI implication: does tightness of some inequalities imply tightness of others? Study the *face lattice* of polyhedral cones spanned by valid linear information inequalities.

## Semimatroids

The (balanced) Shannon information inequalities define a rational polyhedral cone in  $\mathbb{R}^{2^n}$  whose elements are valid inequalities at the Gaussian multiinformation region. The CI inference rules encoded in its face lattice define *semimatroids*.

For  $n = 4$ :

- Gaussoids: 679
- Semimatroidal gaussoids: **629**

For  $n = 5$ :

- Gaussoids: 60,212,776
- Semimatroidal gaussoids: 39,807,192
- Semimatroidal orientable gaussoids: 20,576,142,  
8,148 less than just orientable gaussoids.

**Tools:** Exact rational LP solver to find an inner point on a face of the cone or, alternatively, a Farkas certificate.

# Selfadhesivity

**Theorem.** Positive definite matrices are *selfadhesive* in the sense that for every  $\Sigma \in \text{PD}_N$  and every  $I \subseteq N$  there exist a copy  $M$  of  $N$  with  $N \cap M = I$  and  $\Phi \in \text{PD}_{NM}$  such that:

- $\Phi_N = \Sigma$  and  $\Phi_M = \Sigma$ ,
- $\Phi_{N,M}$  has rank exactly  $|N \cap M|$ .

Such a  $\Phi$  is an *adhesive extension of  $\Sigma$  at  $I$* .

Realizable CI structures inherit *structural* selfadhesivity properties, which are necessary for realizability:

- can be glued as gaussoids (**apparently trivial?**),
- can be glued as orientable gaussoids (**19,723,980/20,584,290**),
- can be glued as semimatroids (**39,595,332/39,807,192**),
- ...

## Outlook: Formal realizations

**Transfer principle.** Let  $\mathbb{K}$  be an ordered field extension of  $\mathbb{R}$ . Then for every matrix  $\Sigma$  over  $\mathbb{K}$  which satisfies

$$\det \Sigma_I > 0 \text{ in } \mathbb{K} \text{ for all } I \subseteq [n],$$

the formal CI structure  $[\Sigma]$  is realizable over  $\mathbb{R}$ .

$$\begin{pmatrix} 1 & g & \varepsilon f & \varepsilon \\ g & 1 & 0 & \varepsilon g \\ \varepsilon f & 0 & 1 & f \\ \varepsilon & \varepsilon g & f & 1 \end{pmatrix}, \quad f = \frac{2\varepsilon}{1 + \varepsilon^2}, \quad g = \frac{1 - \varepsilon^2}{1 + \varepsilon^2}$$

over  $\mathbb{R}(\varepsilon)$ ,  $\varepsilon > 0$  infinitesimal, proves realizability of

$$\{(13|4), (14|23), (23|), (24|1)\}.$$

## Outlook: Optimization

Real algebraic geometry has an exact algebraic solution for the implication / realizability problem (*the real deal, not a relaxation!*):

**Theorem.** The inference rule  $\bigwedge \mathcal{L} \Rightarrow \bigvee \mathcal{M}$  is valid for regular Gaussians if and only if  $\prod_{(ij|K) \in \mathcal{M}} a_{ij|K}$  vanishes on the semialgebraic set of matrices  $\Sigma$  given by





$$\{\det \Sigma_{ij|K} = 0 \ \forall (ij|K) \in \mathcal{L}\} \cap \{\det \Sigma_I > 0 \ \forall I \subseteq [n]\}.$$

This happens if and only if

$$\prod_{(ij|K) \in \mathcal{M}} a_{ij|K} \in \sqrt{\mathcal{I}(a_{ij|K} : (ij|K) \in \mathcal{L}) + \mathcal{I}(1 - y_I^2 p_I : I \subseteq [n])}$$

in the polynomial ring  $\mathbb{R}[\Sigma, y_I : I \subseteq [n]]$ .

**Tools:** Polynomial optimization and SDP relaxations.

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