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► A linear structural equation model defines random variables X recursively via a directed acyclic graph G = (V, E) and Gaussian noise:

$$X_j = \sum_{i \in \mathrm{pa}(j)} \lambda_{ij} X_i + \varepsilon_j, \quad \varepsilon_j \sim \mathcal{N}(0, \omega_j).$$



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- ▶ Parents of node *j* are regarded as direct causes of *j*.
- The vector X is again Gaussian with mean zero. Since G is acyclic, we can solve for the covariance matrix Σ:

$$\Sigma = (I - \Lambda)^{-\mathsf{T}} \Omega (I - \Lambda)^{-1}, \quad \text{with } \Lambda \in \mathbb{R}^{\mathsf{E}} \text{ and } \Omega = \operatorname{diag}(\omega).$$

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▶ All such matrices form the model  $\mathcal{M}(G) \subseteq \mathrm{PD}_V$ .

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- ▶ Model equivalence  $\mathcal{M}(G) = \mathcal{M}(H)$  is combinatorially characterized: if and only if G and H have the same skeleton and v-structures.
  - ► Markov equivalence = ambiguity about the direction of causality.

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- Vertex-only colorings correspond to partial homoscedasticity [WD23].

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► Not invariant anymore.

▶ It follows from the recursive factorization and some linear algebra that

$$\omega_j = \operatorname{Var}(X_i \mid X_{\operatorname{pa}(i)}) = \frac{|\Sigma_{j \cup \operatorname{pa}(j)}|}{|\Sigma_{\operatorname{pa}(j)}|}$$

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► Study the rational functions

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► A set A is identifying for a vertex j resp. edge ij if

$$\omega_j = \omega_{j|A}(\Sigma)$$
 resp.  $\lambda_{ij} = \lambda_{ij|A}(\Sigma)$ 

for all  $\Sigma \in \mathcal{M}(G)$ .

#### Theorem

Let G = (V, E) be a DAG. Then:

•  $\omega_j = \omega_{j|A}(\Sigma)$  for every  $\Sigma \in \mathcal{M}(G)$  if and only if  $pa(j) \subseteq A \subseteq V \setminus \overline{de}(j)$ . [WD23]

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► The polynomials  $\operatorname{vcr}(i|A, j|B) = |\Sigma_A||\Sigma_B|(\omega_{i|A} - \omega_{j|B})$  resp.  $\operatorname{ecr}(ij|A, kl|B) = |\Sigma_A||\Sigma_B|(\lambda_{ij|A} - \lambda_{kl|B})$  vanish on the model  $\mathcal{M}(G, c)$  whenever c(i) = c(j) resp. c(ij) = c(kl) and A and B are identifying.

### **Model geometry**

#### Theorem

For every colored DAG (G, c) the model  $\mathcal{M}(G, c)$  has irreducible Zariski closure and is a smooth submanifold of  $PD_V$ . It is diffeomorphic to an open ball of dimension vc + ec (the number of vertex- and edge-color classes).

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The vanishing ideal  $P_{G,c}$  of  $\mathcal{M}(G,c)$  is  $(I_G + I_c) : S_G$  where:

- ►  $I_G = \langle |\Sigma_{ij|pa(j)}| : ij \notin E \rangle$  is the conditional independence ideal of G,
- ►  $I_c = \langle vcr(i|pa(i), j|pa(j)) : c(i) = c(j) \rangle + \langle ecr(ij|pa(j), kl|pa(l)) : c(ij) = c(kl) \rangle$ is the coloring ideal of G,
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- $S_G = \{\prod_{j \in V} |\Sigma_{pa(j)}|^{k_j} : k_j \in \mathbb{N}\}$  is the monoid of parental principal minors.

Resolves the colored generalization of a conjecture of Sullivant; see also [RP14].

#### Lemma

Let R, R' be rings,  $S \subseteq R$  multiplicatively closed, and:

▶ maps  $\phi : R \to R'$  and  $\psi : R' \to S^{-1}R$  with  $\psi \circ \phi = id_R$ ,

► for a prime ideal  $I' = \langle f_1, \ldots, f_k \rangle$ , write  $\psi(f_i) = g_i/h_i$  and set  $J = \langle g_i \rangle$ . If  $I := \phi^{-1}(I') \in \text{Spec}(S^{-1}R/J)$ , then I = J : S.

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- ► The lemma computes the vanishing ideal up to a saturation of rationally identifiable models with additional equation constraints.
- Knowing a parametrization and generators for the vanishing ideal up to saturation is sufficient in practice for model distinguishability.
- ► Conceivable to extend to non-linear equations and inequalities.

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#### Theorem ([WD23; STD10])



• Generic  $\Sigma \in \mathcal{M}(G, c)$  is faithful to c.

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► The example on the right colors vertices and edges. The generic matrix in the model satisfies 1 ⊥⊥ 4 | 5. No faithful distribution!



#### Theorem ([WD23])

If (G, c) and (H, c) are vertex-colored DAGs, then  $\mathcal{M}(G, c) = \mathcal{M}(H, c)$  if and only if G and H are Markov-equivalent and  $pa_G(j) = pa_H(j)$  for all  $j \in V$  with  $|c^{-1}(j)| \ge 2$ .

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An edge-colored DAG (G, c) is BPEC if:

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- ▶ blocked: color classes partition parent sets of nodes.

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#### Theorem

If (G, c) and (H, d) are two BPEC-DAGs, then  $\mathcal{M}(G, c) = \mathcal{M}(H, d)$  implies (G, c) = (H, d). In particular, the Markov-equivalence classes of BPEC-DAGs are singletons and the causal structure is identifiable.

#### References

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