

WISST! The Lefschetz principle

or: What is... quantifier elimination?

Given polynomials f_1, \dots, f_r and $g_1, \dots, g_s \in \mathbb{Z}[x_1, \dots, x_n]$
study the solvability of the system

$$\begin{aligned} f_i &= 0, \quad i \in [r], \\ g_j &\neq 0, \quad j \in [s], \end{aligned} \tag{\exists}$$

over varying fields \mathbb{K} .

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Note: the \mathbb{Z} coefficients make sense in every field \mathbb{K} by interpreting $k \in \mathbb{Z}$ as $\underbrace{1 + 1 + \dots + 1}_{k \text{ times}} \in \mathbb{K}$.

The Lefschetz principle

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3. In particular if (\exists) has a solution over \mathbb{C} , it has a solution in finite fields \mathbb{F}_{p^m} for all but finitely many primes p .
4. The set of characteristics over which (\exists) has a solution can be effectively computed from the defining polynomials f_i and g_j .

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(First-order) language of rings:

- ▶ constants 0 and 1
- ▶ functions $+$, $-$, \cdot
- ▶ relations $=$
- ▶ Boolean logic connectives \wedge , \vee , \neg (\Rightarrow , \Leftrightarrow , \dots)
- ▶ \exists and \forall quantifiers
- ▶ variables x_1, x_2, \dots

Formulas, axioms and definability

The definition of a ring can be written in the language of rings:

- ▶ $\forall a : a - a = 0$
- ▶ $\forall a, b, c : (a + b) + c = a + (b + c)$
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(\exists) is expressible as a sentence in this language:

$$\exists x_1, \dots, x_n : \bigwedge_{i=1}^r f_i(x_1, \dots, x_n) = 0 \wedge \bigwedge_{j=1}^s \neg(g_j(x_1, \dots, x_n) = 0)$$

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Eliminating quantifiers (and hence **all variables**) from (\exists) results in a Boolean combination of (in)equations $n = m$ for some $n, m \in \mathbb{Z}$.

These inequalities point out exactly which **characteristics** are required and which are ruled out for having a solution to (\exists) .