

The Gaussian CI inference problem

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Non-linear algebra seminar

14 October 2021

Gaussian conditional independence

Consider random variables $(\xi_i)_{i \in N}$. The *conditional independence (CI) statement* $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ conveys, informally, that if ξ_K is known, then learning the value of ξ_i does not give any information about ξ_j .

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Example: Let c_1 and c_2 be two independent coins and b a bell which rings if and only if c_1 and c_2 land with the same side up. What is the conditional independence relation of the system (c_1, c_2, b) of random variables?

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Example: Let c_1 and c_2 be two independent coins and b a bell which rings if and only if c_1 and c_2 land with the same side up. What is the conditional independence relation of the system (c_1, c_2, b) of random variables? $\rightarrow c_1 \perp\!\!\!\perp c_2$ and $\neg(c_1 \perp\!\!\!\perp c_2 \mid b)$...

Gaussian conditional independence

Let the random vector be normally distributed with covariance matrix $\Sigma \in \text{PD}_N$.

Definition

The polynomial $\Sigma[K] := \det \Sigma_{K,K}$ is a *principal minor* of Σ and $\Sigma[ij|K] := \det \Sigma_{iK,jK}$ is an *almost-principal minor*.

- ▶ Σ is PD if and only if $\Sigma[K] > 0$ for all $K \subseteq N$.
- ▶ $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ holds if and only if $\Sigma[ij|K] = 0$.

Special polynomials

$$\Sigma[ij|] = x_{ij}$$

$$\Sigma[ij|k] = x_{ij}x_{kk} - x_{ik}x_{jk}$$

$$\Sigma[ij|kl] = x_{ij}x_{kk}x_{ll} - x_{il}x_{jl}x_{kk} + x_{il}x_{jk}x_{kl} + x_{ik}x_{jl}x_{kl} - x_{ij}x_{kl}^2 - x_{ik}x_{jk}x_{ll}$$

$$\begin{aligned}\Sigma[ij|klm] = & x_{ij}x_{kk}x_{ll}x_{mm} + x_{im}x_{jm}x_{kl}^2 - x_{im}x_{jl}x_{kl}x_{km} - x_{il}x_{jm}x_{kl}x_{km} + x_{il}x_{jl}x_{km}^2 \\ & - x_{im}x_{jm}x_{kk}x_{ll} + x_{im}x_{jk}x_{km}x_{ll} + x_{ik}x_{jm}x_{km}x_{ll} - x_{ij}x_{km}^2x_{ll} \\ & + x_{im}x_{jl}x_{kk}x_{lm} + x_{il}x_{jm}x_{kk}x_{lm} - x_{im}x_{jk}x_{kl}x_{lm} - x_{ik}x_{jm}x_{kl}x_{lm} \\ & - x_{il}x_{jk}x_{km}x_{lm} - x_{ik}x_{jl}x_{km}x_{lm} + 2x_{ij}x_{kl}x_{km}x_{lm} + x_{ik}x_{jk}x_{lm}^2 \\ & - x_{ij}x_{kk}x_{lm}^2 - x_{il}x_{jl}x_{kk}x_{mm} + x_{il}x_{jk}x_{kl}x_{mm} + x_{ik}x_{jl}x_{kl}x_{mm} \\ & - x_{ij}x_{kl}^2x_{mm} - x_{ik}x_{jk}x_{ll}x_{mm}\end{aligned}$$

⋮

Gaussian CI models

Definition

A *CI constraint* is a CI statement $\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K$ or its negation $\neg(\xi_i \perp\!\!\!\perp \xi_j \mid \xi_K)$.

The *model* of a set of CI constraints is the set of all PD matrices which satisfy them.

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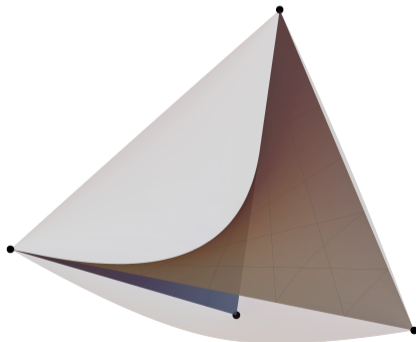


Figure: Model of $\Sigma[12|3] = 0$ in the space of 3×3 correlation matrices.

Models and inference

Consider two sets of CI statements \mathcal{P} and \mathcal{Q} :

$$\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$$

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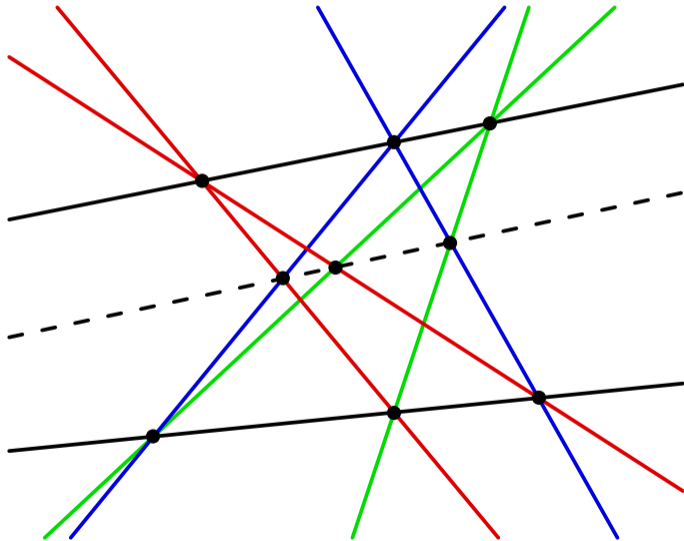
Models and inference

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Reasoning about relevance statements in normally distributed random variables is **the same** as reasoning about the vanishing of very special kinds of determinants on very special kinds of varieties inside the positive-definite matrices.

For geometers: conditional independence \approx collinearity



Examples of CI inference

Consider a general positive-definite 3×3 correlation matrix

$$\Sigma = \begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix}.$$

- ▶ If $\Sigma[12|3] = a - bc$ and $\Sigma[13|] = b$ vanish, then $\Sigma[12|] = a$ and $\Sigma[13|2] = b - ac$ must vanish as well:

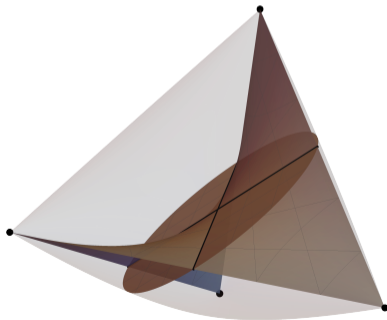
$$[12|3] \wedge [13|] \Rightarrow [12|] \wedge [13|2].$$

Examples of CI inference

$$\Sigma = \begin{pmatrix} 1 & a & b \\ a & 1 & c \\ b & c & 1 \end{pmatrix}$$

- ▶ If $\Sigma[12|] = a$ and $\Sigma[12|3] = a - bc$ vanish, then $bc = \Sigma[13|] \cdot \Sigma[23|]$ must vanish:

$$[12|] \wedge [12|3] \Rightarrow [13|] \vee [23|].$$



Rational points on CI models

Šimeček's Question (2006)

Does every non-empty Gaussian CI model contain a rational point?

Or: can every wrong inference rule be refuted over \mathbb{Q} ?

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Model M85



1 a b c
a 1 d e
b d 1 f
c e f 1

Where:

$$a = \frac{3}{632836} \sqrt{1107463},$$

$$b = 10c = \frac{100}{158209} \sqrt{1107463}$$

$$d = 10e = \frac{3}{4}, f = \frac{1}{10}$$

$$\begin{pmatrix} 1 & -1/17 & -49/51 & -7/17 \\ -1/17 & 1 & 1/3 & 1/7 \\ -49/51 & 1/3 & 1 & 3/7 \\ -7/17 & 1/7 & 3/7 & 1 \end{pmatrix}$$

Complexity bounds from real geometry

Let $f_i \in \mathbb{Z}[t_1, \dots, t_k]$ be integer polynomials in finitely many variables.

Theorem (Tarski's transfer principle)

If a polynomial system $\{f_i \bowtie_i 0\}$, where $\bowtie_i \in \{=, \neq, <, \leq, \geq, >\}$, has a solution over \mathbb{R} , then it has a solution in a finite real extension of \mathbb{Q} .

→ If $\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$ is false, there exists a counterexample matrix Σ with algebraic entries.

$[12 |] \wedge [12 | 3] \Rightarrow [13 |]$ is false and a counterexample is

$$\begin{pmatrix} 1 & 0 & 1/2 \\ 0 & 1 & 0 \\ 1/2 & 0 & 1 \end{pmatrix}.$$

Complexity bounds from real geometry

Let $f_i, g_j, h_k \in \mathbb{Z}[t_1, \dots, t_k]$ be integer polynomials in finitely many variables.

Theorem (Positivstellensatz)

A polynomial system $\{f_i = 0, g_j \geq 0, h_k \neq 0\}$ is infeasible if and only if there exist $f \in \text{ideal}(f_i)$, $g \in \text{cone}(g_j)$ and $h \in \text{monoid}(h_k)$ such that $g + h^2 = f$.

→ If $\bigwedge \mathcal{P} \Rightarrow \bigvee \mathcal{Q}$ is **true**, there exists an algebraic proof for it with integer coefficients.

$[12|] \wedge [12|3] \Rightarrow [13|] \vee [23|]$ is true and a proof is the polynomial identity

$$\Sigma[13|] \cdot \Sigma[23|] = \Sigma[3] \cdot \Sigma[12|] - \Sigma[12|3].$$

The associated decision problem is the **existential theory of the reals**.

A 5×5 final polynomial

The following inference rule is valid for all positive-definite 5×5 matrices:

$$[12|] \wedge [14|5] \wedge [23|5] \wedge [35|1] \wedge [45|2] \wedge [15|23] \wedge [34|12] \wedge [24|135] \Rightarrow [25|] \vee [34|].$$

A 5×5 final polynomial

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$$\begin{aligned} & [25 |] [34 |] \cdot [1][2][3][15] = \\ & \left(cd^2egr + bd^2fgr - ad^2grh - 2cd^2e^2i - 2bd^2efi - 2pdfgri + 2ad^2ehi + 2pdef^2 - 2pdqhi^2 + 2pcqi^3 + \right. \\ & \left. 2pdqrij - 2pbqi^2j - pcegrt + pbfgrt + pagrht + 2pce^2it - 2pcqrit + 2pbqhit - 2paehit \right) \cdot [12 |] + \\ & \left(pdqer + pbqgr - 2pbqei \right) \cdot [14 | 5] - \left(pcdqr + p^2fgr - 2pbcqi + 2pb^2qj - 2p^2qrj \right) \cdot [23 | 5] + \\ & \left(cdqgr - 2cdqei + 2pqghi - 2pqfi^2 - pqgrj + 2pqueij - 2pe^2ft + 2pqfrt \right) \cdot [35 | 1] + \\ & \left(pd^2er - 2pbdei + p^2gri + 2pb^2et - 2p^2ert \right) \cdot [45 | 2] - \left(2pdfi - 2pbft \right) \cdot [15 | 23] - \\ & \left(d^2gr - 2d^2ei - pgrt + 2peit \right) \cdot [34 | 12] - 2pqi \cdot [24 | 135]. \end{aligned}$$

Universality theorems

Theorem (B. 2021)

For every finite real extension \mathbb{K}/\mathbb{Q} there exists a Gaussian CI model $\mathcal{M}_{\mathbb{K}}$ such that: for every \mathbb{L}/\mathbb{Q} , $\mathcal{M}_{\mathbb{K}}$ has an \mathbb{L} -rational point if and only if $\mathbb{K} \subseteq \mathbb{L}$.

→ The answer to Šimeček's question is **NO**.

Theorem (B. 2021)

The problem of deciding whether a CI inference formula is valid for all Gaussian distributions is polynomial-time equivalent to the existential theory of the reals.

Approximations to the inference problem

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Theorem (Matúš 2005)

The following relations hold for every symmetric matrix Σ :

$$\begin{aligned}\Sigma[ij|L]^2 &= \Sigma[iL] \cdot \Sigma[jL] - \Sigma[L] \cdot \Sigma[ijL] \\ \Sigma[kL] \cdot \Sigma[ij|L] &= \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]\end{aligned}$$

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$$\begin{aligned}\Sigma[ij|L]^2 &= \Sigma[iL] \cdot \Sigma[jL] - \Sigma[L] \cdot \Sigma[ijL] \quad \rightarrow \text{semimatroids} \\ \Sigma[kL] \cdot \Sigma[ij|L] &= \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L] \quad \rightarrow \text{gaussoids}\end{aligned}$$

These relations define essential geometric properties of symmetric matrices in principal and almost-principal minor coordinates. Study their combinatorics!

The Gaussian CI configuration space

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

The *Gaussian CI configuration space* $\mathcal{G} \subseteq \mathbb{R}^{2^n} \times \mathbb{R}^{\binom{n}{2}2^{n-2}}$ consists of all vectors of principal and almost-principal minors of $\Sigma \in \text{PD}_n$.

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Very wasteful encoding of a matrix, but this creates simple and useful relations on configuration vectors. The CI structure of Σ is encoded in its *zero pattern*.

Combinatorial compatibility

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

Combinatorial compatibility means **fulfilling polynomial relations under uncertainty**:
What if we only knew that all $\Sigma[K] \neq 0$ and whether or not $\Sigma[ij|K] = 0$?

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$$[ij|L] \wedge [ij|kL] \Rightarrow [ik|L] \vee [jk|L]$$

$$[ik|L] \wedge [ij|kL] \Rightarrow [ij|L]$$

⋮

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$$[ij|L] \wedge [ik|L] \Rightarrow [ij|kL] \wedge [ik|jL]$$

This yields the definition of *gaussoids*.

CI inference via SAT solvers

Since gaussoids have a finite axiomatization, a SAT solver like CaDiCaL can deduce implications under the gaussoid axioms:

$$\begin{aligned} & [12|3] \wedge [12|34] \wedge [24|1] \wedge [34|2] \\ \Rightarrow & [12|] \wedge [12|4] \wedge [24|] \wedge [24|3] \wedge [24|13] \wedge [34|] \end{aligned}$$

These conclusions are valid for all regular Gaussian distributions.

Oriented gaussoids

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

What if we only knew that all $\text{sgn } \Sigma[K] = +1$ and the value of $\text{sgn } \Sigma[ij|K]$?

$$[ij|L] > 0 \wedge [ij|kL] < 0 \Rightarrow ([ik|L] > 0 \wedge [jk|L] > 0) \vee ([ik|L] < 0 \wedge [jk|L] < 0)$$

→ *Oriented and orientable gaussoids.*

Oriented gaussoids

$$\Sigma[kL] \cdot \Sigma[ij|L] = \Sigma[L] \cdot \Sigma[ij|kL] + \Sigma[ik|L] \cdot \Sigma[jk|L]$$

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$$[ij|L] > 0 \wedge [ij|kL] < 0 \Rightarrow ([ik|L] > 0 \wedge [jk|L] > 0) \vee ([ik|L] < 0 \wedge [jk|L] < 0)$$

→ *Oriented* and *orientable* gaussoids.

$$[ij|L] \wedge [kl|L] \wedge [ik|jL] \wedge [jl|ikL] \Rightarrow [ik|L]$$

$$[ij|L] \wedge [kl|iL] \wedge [kl|jL] \wedge [ij|kL] \Rightarrow [kl|L]$$

$$[ij|L] \wedge [jl|kL] \wedge [kl|iL] \wedge [ik|jL] \Rightarrow [ik|L]$$

$$[ij|kL] \wedge [ik|lL] \wedge [il|jL] \Rightarrow [ij|L]$$

$$[ij|kL] \wedge [ik|lL] \wedge [jl|iL] \wedge [kl|jL] \Rightarrow [ij|L]$$

CI inference via SAT solvers II

Using the gaussoid axioms, we find:

$$[12 |] \wedge [13 | 4] \wedge [14 | 5] \wedge [15 | 23] \wedge [23 | 5] \wedge [24 | 135] \wedge [34 | 12] \wedge [35 | 1] \wedge [45 | 2] \\ \Rightarrow \text{nothing.}$$

The structure on the left is a gaussoid.

CI inference via SAT solvers II

Running the SAT solver CaDiCaL on the definition of **oriented** gaussoids confirms that their supports satisfy

$$[12|] \wedge [13|4] \wedge [14|5] \wedge [15|23] \wedge [23|5] \wedge [24|135] \wedge [34|12] \wedge [35|1] \wedge [45|2]$$

\Rightarrow everything except $[25|K]$ for all K .

The geometric model is a Gaussian graphical model!

The search for inference rules (since at least 2008!)

Inference rules help characterize the *realizable* CI structures:

- ▶ 3-variate: 11 out of 64 by Matúš 2005.
- ▶ 4-variate: 629 out of 16 777 216 by Lněnička and Matúš 2007.
- ▶ 5-variate: *open!* (out of 1 208 925 819 614 629 174 706 176)
 - ▶ 254 826 gaussoids modulo symmetry
 - ▶ 87 834 of which are orientable gaussoids
 - ▶ 84 908 of which are [selfadhesive](#) orientable gaussoids.
 - ▶ 84 434 of which are selfadhesive (orientable gaussoids \cap [semimatroids](#)).




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Help wanted:

- ▶ Use information inequalities and linear programming.
- ▶ Tropical approximations and valuated gaussoids.
- ▶ Compute algebraic realization spaces.
- ▶ Find and certify real solutions to polynomial systems.

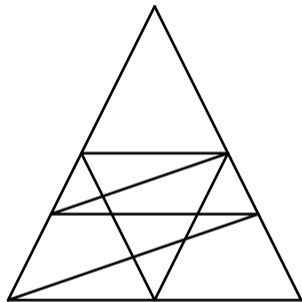
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Proof sketch (1): Algebra \subseteq Synthetic geometry

Point and line configuration for the equation $x^2 - 2 = 0$.

The configuration is specified by incidences between points and lines and also the parallelities of lines.

It is realizable over $\mathbb{Q}(\sqrt{2})$ but not over \mathbb{Q} .



Keyword for the general technique: *von Staudt constructions* (1857).

Proof sketch (2): Synthetic geometry \subseteq Gaussian CI

$\Sigma[ij] = x_{ij} \rightarrow$ impose $x_{kl} = x_{km} = x_{lm} = 0$ on a correlation matrix, then:

$$\begin{aligned}
 \Sigma[ij | klm] &= x_{ij}x_{kk}x_{ll}x_{mm} + x_{im}x_{jm}\underline{x_{kl}^2} - x_{im}x_{jl}\underline{x_{kl}x_{km}} - x_{il}x_{jm}\underline{x_{kl}x_{km}} + x_{il}x_{jl}\underline{x_{km}^2} \\
 &\quad - x_{im}x_{jm}x_{kk}x_{ll} + x_{im}x_{jk}\underline{x_{km}x_{ll}} + x_{ik}x_{jm}\underline{x_{km}x_{ll}} - x_{ij}\underline{x_{km}^2}x_{ll} \\
 &\quad + x_{im}x_{jl}x_{kk}x_{lm} + x_{il}x_{jm}x_{kk}x_{lm} - x_{im}x_{jk}\underline{x_{kl}x_{lm}} - x_{ik}x_{jm}\underline{x_{kl}x_{lm}} \\
 &\quad - x_{il}x_{jk}\underline{x_{km}x_{lm}} - x_{ik}x_{jl}\underline{x_{km}x_{lm}} + 2x_{ij}\underline{x_{kl}x_{km}x_{lm}} + x_{ik}x_{jk}\underline{x_{lm}^2} \\
 &\quad - x_{ij}x_{kk}\underline{x_{lm}^2} - x_{il}x_{jl}x_{kk}x_{mm} + x_{il}x_{jk}\underline{x_{kl}x_{mm}} + x_{ik}x_{jl}\underline{x_{kl}x_{mm}} \\
 &\quad - x_{ij}\underline{x_{kl}^2}x_{mm} - x_{ik}x_{jk}x_{ll}x_{mm} \\
 &= x_{ij} - \sum_{t=k,l,m} x_{it}x_{jt} = x_{ij} - \left\langle \begin{pmatrix} x_{ik} \\ x_{il} \\ x_{im} \end{pmatrix}, \begin{pmatrix} x_{jk} \\ x_{jl} \\ x_{jm} \end{pmatrix} \right\rangle.
 \end{aligned}$$

Incidence relation in a CI model

$$\begin{array}{c}
 p_1 \\
 \vdots \\
 p_n \\
 l_1 \\
 \vdots \\
 l_m \\
 x \\
 y \\
 z
 \end{array}
 \left(
 \begin{array}{ccc|ccc|ccc}
 & p_1 & \dots & p_n & l_1 & \dots & l_m & x & y & z \\
 p_1 & p_1^* & & \langle p, p' \rangle & & & & p_1^x & p_1^y & p_1^z \\
 \vdots & & \ddots & & & \langle p, l \rangle & & & \vdots & \\
 p_n & \langle p', p \rangle & & p_n^* & & & & p_n^x & p_n^y & p_n^z \\
 \hline
 l_1 & & & & l_1^* & & \langle l, l' \rangle & l_1^x & l_1^y & l_1^z \\
 \vdots & & \langle l, p \rangle & & & \ddots & & & \vdots & \\
 l_m & & & & \langle l', l \rangle & & l_m^* & l_m^x & l_m^y & l_m^z \\
 \hline
 x & p_1^x & & p_n^x & l_1^x & & l_m^x & 1 & 0 & 0 \\
 y & p_1^y & \dots & p_n^y & l_1^y & \dots & l_m^y & 0 & 1 & 0 \\
 z & p_1^z & & p_n^z & l_1^z & & l_m^z & 0 & 0 & 1
 \end{array}
 \right)$$